

# The Complexity of Counting Eulerian Tours in 4-Regular Graphs<sup>\*</sup>

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**Abstract.** We investigate the complexity of counting Eulerian tours (#ET) and its variations from two perspectives—the complexity of exact counting and the complexity w.r.t. approximation-preserving reductions (AP-reductions [DGGJ04]). We prove that #ET is #P-complete even for planar 4-regular graphs.

A closely related problem is that of counting A-trails (#A-TRAILS) in graphs with rotational embedding schemes (so called maps). Kotzig [Kot68] showed that #A-TRAILS can be computed in polynomial time for 4-regular plane graphs (embedding in the plane is equivalent to giving a rotational embedding scheme). We show that for 4-regular maps the problem is #P-hard. Moreover, we show that from the approximation viewpoint #A-TRAILS in 4-regular maps captures the essence of #ET, that is, we give an AP-reduction from #ET in general graphs to #A-TRAILS in 4-regular maps. The reduction uses a fast mixing result for a card shuffling problem [Wil04].

## 1 Introduction

An Eulerian tour in a graph is a tour which travels each edge exactly once. The problem of counting Eulerian tours (#ET) of a graph is one of a few recognized counting problems (see, e.g., [Vaz01], p. 339). The exact counting in general graphs is #P-complete [BW05], and thus there is no polynomial-time algorithm for it unless  $P=NP$ . For the approximate counting one wants to have a fully polynomial randomized approximation scheme (FPRAS), that is, an algorithm which on every instance  $x$  of the problem and error parameter  $\varepsilon > 0$ , will output a value within a factor  $\exp(\pm\varepsilon)$  of  $f(x)$  with probability at least  $2/3$  and in time polynomial in the length of the encoding of  $x$  and  $1/\varepsilon$ , where  $f(x)$  is the value we want to compute. The existence of an FPRAS for #ET is an open problem [TV01, Jer, Vaz01].

A closely related problem to #ET is the problem of counting A-trails (#A-TRAILS) in graphs with rotational embedding schemes (called maps, see Section 2 for a definition). A-trails were studied in the context of decision problems (for example, it is NP-complete to decide whether a given plane graph has an A-trail [BM87, AF95]; on the other hand for 4-regular maps the problem is in

P [Dvo04]), as well as counting problems (for example, Kotzig [Kot68] showed that #A-TRAILS can be computed in polynomial time for 4-regular plane graphs, reducing the problem to counting of spanning trees).

In this paper, we investigate the complexity of #ET in 4-regular graphs and its variations from two perspectives. First, the complexity of exact counting is considered. We prove that #ET in 4-regular graphs (even in 4-regular planar graphs) is #P-complete. We also prove that #A-TRAILS in 4-regular maps is #P-complete (recall that the problem can be solved in polynomial time for 4-regular *plane* graphs).

The second perspective is the complexity w.r.t. the AP-reductions proposed by Dyer, Goldberg, Greenhill and Jerrum [DGGJ04]. We give an AP-reduction from #ET in general graphs to #A-TRAILS in 4-regular maps. Thus we show that if there is an FPRAS for #A-TRAILS in 4-regular maps, then there is also an FPRAS for #ET in general graphs. The existence of AP-reduction from #ET in general graphs to #ET in 4-regular graphs is left open.

## 2 Definitions and Terminology

For the definitions of cyclic orderings, A-trails, and mixed graphs, we follow [Fle90]. Let  $G = (V, E)$  be a graph. For a vertex  $v \in V$  of degree  $d > 0$ , let  $K(v) = \{e_1, \dots, e_d\}$  be the set of edges adjacent to  $v$  in  $G$ . The *cyclic ordering*  $O^+(v)$  of the edges adjacent to  $v$  is a  $d$ -tuple  $(e_{\sigma(1)}, \dots, e_{\sigma(d)})$ , where  $\sigma$  is a permutation in  $S_d$ . We say  $e_{\sigma(i)}$  and  $e_{\sigma(i+1)}$  are cyclicly-adjacent in  $O^+(v)$ , for  $1 \leq i \leq d$ , where we set  $\sigma(d+1) := \sigma(1)$ . The set  $O^+(G) = \{O^+(v) | v \in V\}$  is called a *rotational embedding scheme* of  $G$ . For a plane graph  $G = (V, E)$ , if  $O^+(v)$  is not specified, we usually set  $O^+(v)$  to be the clockwise order of the half-edges adjacent to  $v$  for each  $v \in V$ .

Let  $G = (V, E)$  be a graph with a rotational embedding scheme  $O^+(G)$ . An Eulerian tour  $v_0, e_1, v_1, e_2, \dots, e_\ell, v_\ell = v_0$  is called an *A-trail* if  $e_i$  and  $e_{i+1}$  are cyclicly-adjacent in  $O^+(v_i)$ , for each  $1 \leq i \leq \ell$ , where we set  $e_{\ell+1} := e_1$ .

Let  $G = (V, E, E')$  be a mixed graph, that is,  $E$  is the set of edges and  $E'$  is the set of half-edges (which are incident with only one vertex in  $V$ ). Let  $|E'| = 2d$  where  $d$  is a positive integer and assume that the half-edges in  $E'$  are labelled by numbers from 1 to  $2d$ . A *route*  $r(a, b)$  is a trail (no repeated edges, repeated vertices allowed) in  $G$  that starts with half-edge  $a$  and ends with half-edge  $b$ . A collection of  $d$  routes is called *valid* if every edge and every half-edge is travelled exactly once.

We say that a valid set of routes is of the *type*  $\{\{a_1, b_1\}, \dots, \{a_d, b_d\}\}$  if it contains routes connecting  $a_i$  to  $b_i$  for  $i \in [d]$ . We use  $VR(\{\{a_1, b_1\}, \dots, \{a_d, b_d\}\})$  to denote the set of valid sets of routes of type  $\{\{a_1, b_1\}, \dots, \{a_d, b_d\}\}$  in  $G$ .

We will use the following concepts from Markov chains to construct the gadget in Section 4 (see, e. g., [Jer03] for more detail). Given two probability distributions  $\pi$  and  $\pi'$  on finite set  $\Omega$ , the *total variation distance* between  $\pi$  and  $\pi'$  is defined as

$$\|\pi - \pi'\|_{TV} = \frac{1}{2} \sum_{\omega \in \Omega} |\pi(\omega) - \pi'(\omega)| = \max_{A \subseteq \Omega} |\pi(A) - \pi'(A)|.$$

Given a finite ergodic Markov chain with transition matrix  $P$  and stationary distribution  $\pi$ , the *mixing time* from initial state  $x$ , denoted as  $\tau_x(\varepsilon)$ , is defined as

$$\tau_x(\varepsilon) = \min\{t : \|P^t(x, \cdot) - \pi\|_{TV} \leq \varepsilon\},$$

and the mixing time of the chain  $\tau(\varepsilon)$  is defined as

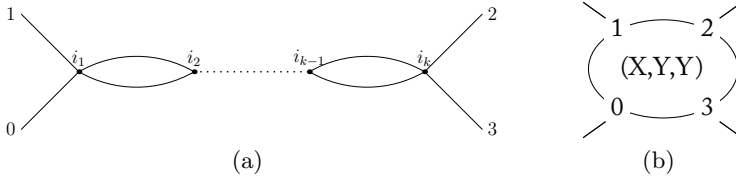
$$\tau(\varepsilon) = \max_{x \in \Omega} \{\tau_x(\varepsilon)\}.$$

### 3 The Complexity of Exact Counting

#### 3.1 Basic Gadgets

We describe two basic gadgets and their properties which will be used as a basis for larger gadgets in the subsequent sections.

The first gadget, which is called the  $(X, Y, Y)$  node, is shown in Figure 1(a), and it is represented by the symbol shown in Figure 1(b). There are  $k$  internal vertices in the gadget, and the labels 0, 1, 2 and 3 are four half-edges of the  $(X, Y, Y)$  node which are the only connections from the outside.



**Fig. 1.** An  $(X, Y, Y)$  node and its symbol. (a): an  $(X, Y, Y)$  node consisting of  $k$  internal vertices; (b): symbol representing the  $(X, Y, Y)$  node.

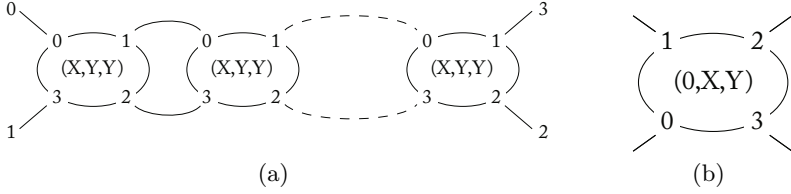
By elementary counting we obtain the following fact.

**Lemma 1.** *The  $(X, Y, Y)$  node with parameter  $k$  has three different types of valid sets of routes and these satisfy*

$$\begin{aligned} |VR(\{0, 1\}, \{2, 3\})| &= k2^{k-1}, \\ |VR(\{0, 2\}, \{1, 3\})| &= |VR(\{0, 3\}, \{1, 2\})| = 2^{k-1}. \end{aligned}$$

*The gadget has  $k$  vertices.*

The second gadget, which is called the  $(0, X, Y)$  node, is shown in Figure 2(a), and it is represented by the symbol shown in Figure 2(b). Let  $p$  be any odd prime. In the construction of the  $(0, X, Y)$  node we use  $p$  copies of  $(X, Y, Y)$



**Fig. 2.** A  $(0, X, Y)$  node and its symbol. (a): a  $(0, X, Y)$  node consisting of  $p$  copies of  $(X, Y, Y)$  nodes; (b): symbol representing the  $(0, X, Y)$  node.

nodes as basic components, and each  $(X, Y, Y)$  node has the same parameter  $k$ . As illustrated, half-edges are connected between two consecutive  $(X, Y, Y)$  nodes. The four labels 0, 1, 2 and 3 at four corners in Figure 2(a) are the four half-edges of the  $(0, X, Y)$  node, and they are the only connections from the outside.

By elementary counting, binomial expansion, Fermat's little theorem, and the fact that 2 has a multiplicative inverse mod  $p$ , we obtain the following:

**Lemma 2.** *Let  $p$  be an odd prime and let  $k$  be an integer. The  $(0, X, Y)$  node with parameters  $p$  and  $k$  has three different types of valid sets of routes and these satisfy*

$$|VR(\{0, 1\}, \{2, 3\})| = pA(A + B)^{p-1} \equiv 0 \pmod{p}, \quad (1)$$

$$|VR(\{0, 2\}, \{1, 3\})| = \frac{(A + B)^p - (B - A)^p}{2} \equiv A \pmod{p}, \quad (2)$$

$$|VR(\{0, 3\}, \{1, 2\})| = \frac{(A + B)^p + (B - A)^p}{2} \equiv B \pmod{p}, \quad (3)$$

where  $A = 2^{k-1}$  and  $B = k2^{k-1}$ . The gadget has  $kp$  vertices.

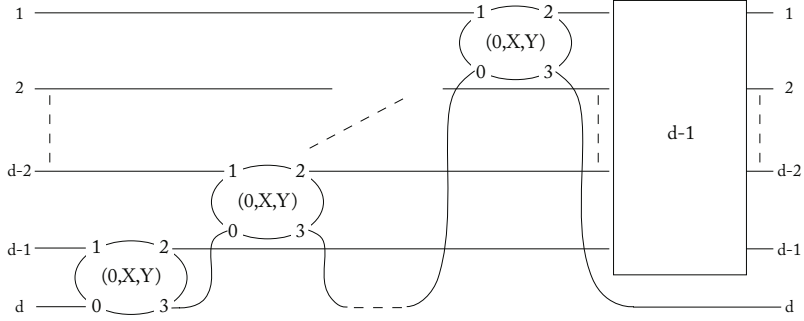
### 3.2 #ET in 4-Regular Graphs Is #P-Complete

Next, we will give a reduction from #ET in general Eulerian graphs to #ET in 4-regular graphs.

**Theorem 1.** *#ET in general Eulerian graphs is polynomial time Turing reducible to #ET in 4-regular graphs.*

The proof of Theorem 1 is postponed to the end of this section.

We use the gadget, which we will call  $Q$ , illustrated in Figure 3 to prove the Theorem. The gadget is constructed in a recursive way. The  $d$  labels  $1, \dots, d$  on the left are called input half-edges of the gadget, and the  $d$  labels on the right are called output half-edges. Given a prime  $p$  and a positive integer  $d$ , the gadget consists of  $d - 1$  copies of  $(0, X, Y)$  nodes with different parameters and one recursive part represented by a rectangle with  $d - 1$  input half-edges and  $d - 1$  output ones. For  $1 \leq i \leq d - 1$ , the  $i$ -th  $(0, X, Y)$  node from left has parameters



**Fig. 3.** Gadget  $Q$  with  $d$  input half-edges and  $d$  output half-edges

$p$  and  $i$ . Half-edge 0 of the  $i$ -th  $(0, X, Y)$  node is connected to half-edge 3 of the  $(i - 1)$ -st  $(0, X, Y)$  node except that for the 1st  $(0, X, Y)$  node half-edge 0 is the  $d$ -th input half-edge of the gadget. Half-edge 1 of the  $i$ -th  $(0, X, Y)$  node is the  $(d - i)$ -th input half-edge of the gadget. Half-edge 2 of the  $i$ -th  $(0, X, Y)$  node is connected to the  $(d - i)$ -th input half-edge of the rectangle. Half-edge 3 of the  $(d - 1)$ -st  $(0, X, Y)$  node is the  $d$ -th output half-edge of the gadget. For  $1 \leq j \leq d - 1$ , the  $j$ -th output half-edge of the rectangle is the  $j$ -th output half-edge of the gadget. From the constructions of  $(X, Y, Y)$  nodes and  $(0, X, Y)$  nodes, the total size of the  $d - 1$  copies of  $(0, X, Y)$  nodes is  $O(pd^2)$ . Thus, the size of the gadget is  $O(pd^3)$ .

**Lemma 3.** Consider the gadget  $Q$  with parameters  $d$  and  $p$ . Let  $\sigma$  be a permutation in  $S_d$ . Then

$$|VR(\sigma)| := |VR(\{IN_1, OUT_{\sigma(1)}\}, \dots, \{IN_d, OUT_{\sigma(d)}\})| \equiv R_d \pmod{p}, \quad (4)$$

where  $R_d \equiv \prod_{i=1}^{d-1} (2^{i(i-1)/2} i!)$ .

Moreover, any type  $\tau$  which connects two  $IN$  (or two  $OUT$ ) half-edges satisfies

$$|VR(\tau)| \equiv 0 \pmod{p}. \quad (5)$$

*Proof.* The proof is by induction on  $d$ , the base case  $d = 1$  is trivial. Suppose the statement is true for gadget  $Q$  with  $(d - 1)$  input half-edges, that is,  $|VR(\varrho)| \equiv R_{d-1} \pmod{p}$  for every  $\varrho \in S_{d-1}$ .

Now, consider gadget  $Q$  with  $d$  input half-edges. For  $1 \leq j \leq d - 1$ , we cut the gadget by a vertical line just after the  $j$ -th  $(0, X, Y)$  node and only consider the part of the gadget to the left of the line, we will call this partial gadget  $Q_j$ .

*Claim.* Let  $A_s$  be the set of permutations in  $S_d$  which map  $s$  to  $d$ . In the partial gadget  $Q_j$  we have that for  $s \in \{d - j, \dots, d\}$  have

$$\sum_{\sigma \in A_s} |VR_{Q_j}(\sigma)| \equiv j! 2^{j(j-1)/2} \pmod{p},$$

where the subscript  $Q_j$  is used to indicate that we count routes in gadget  $Q_j$ .

*Proof of the Claim.* We prove the claim by induction on  $j$ , the base case  $j = 1$  is trivial.

Now assume that the claim is true for  $j-1$ , that is, for all  $s \in \{d-j+1, \dots, d\}$  in gadget  $Q_{j-1}$  we have

$$\sum_{\sigma \in A_s} |VR_{Q_{j-1}}(\sigma)| \equiv (j-1)!2^{(j-1)(j-2)/2} \pmod{p}.$$

The  $j$ -th  $(0, X, Y)$  node takes  $(d-j)$ -th input half-edge of the gadget and the half-edge 3 of the  $(j-1)$ -st  $(0, X, Y)$  node, and has parameters  $p$  and  $j$ .

The type of the  $j$ -th  $(0, X, Y)$  node is  $\{\{0, 2\}, \{1, 3\}\}$  if and only if the resulting permutation in  $Q_j$  is in  $A_{d-j}$ . Thus we have

$$\sum_{\sigma \in A_{d-j}} |VR_{Q_j}(\sigma)| \equiv 2^{j-1} \prod_{k=1}^{j-1} (2^{k-1}(k+1)) \equiv j!2^{j(j-1)/2} \pmod{p},$$

where the first term is the number of choices (modulo  $p$ ) in the  $j$ -th  $(0, X, Y)$  node to make it  $\{\{0, 2\}, \{1, 3\}\}$  and the  $k$ -th term in the product is the number of choices (modulo  $p$ ) in the  $k$ -th  $(0, X, Y)$  node to make it either  $\{\{0, 2\}, \{1, 3\}\}$  or  $\{\{0, 3\}, \{1, 2\}\}$ .

If the type inside the  $j$ -th  $(0, X, Y)$  node is  $\{\{0, 3\}, \{1, 2\}\}$  then the resulting permutation is in  $A_s$  for  $s \in \{d-j+1, \dots, d\}$ . Thus

$$\begin{aligned} \sum_{\sigma \in A_s} |VR_{Q_j}(\sigma)| &\equiv j2^{j-1} \sum_{\sigma \in A_s} |VR_{Q_{j-1}}(\sigma)| \equiv j2^{j-1}(j-1)!2^{(j-1)(j-2)/2} \\ &\equiv j!2^{j(j-1)/2} \pmod{p}, \end{aligned}$$

where  $j2^{j-1}$  is the number of choices (modulo  $p$ ) in the  $j$ -th  $(0, X, Y)$  node to make it  $\{\{0, 3\}, \{1, 2\}\}$ .  $\square$

Now we continue with the proof of the Lemma 3.

Let  $\sigma$  be a permutation in  $S_d$ . Let  $l = \sigma^{-1}(d)$ . In order for  $\sigma$  to be realized by gadget  $Q$  we have to have  $l$  mapped to  $d$  by  $Q_{d-1}$  and the permutation realized by the recursive gadget of size  $d-1$  must “cancel” the permutation of  $Q_{d-1}$ . By the claim there are  $(d-1)!2^{(d-1)(d-2)/2}$  (modulo  $p$ ) choices in  $Q_{d-1}$  which map  $l$  to  $d$  and by the inductive hypothesis there are  $R_{d-1}$  (modulo  $p$ ) choices in the recursive gadget of size  $d-1$  that give the unique permutation that “cancels” the permutation of  $Q_{d-1}$ . Thus

$$|VR(\sigma)| \equiv R_d \equiv (d-1)!2^{(d-1)(d-2)/2} R_{d-1} \pmod{p},$$

finishing the proof of (4).

To see (5) note that the number of valid sets of routes which contain route starting and ending at both input half-edges or both output half-edges is 0 modulo  $p$ . This is because the number of valid set of routes of type  $\{\{0, 1\}, \{2, 3\}\}$  inside the  $(0, X, Y)$  node is 0 modulo  $p$ .  $\square$

*Proof (of Theorem 1).* The reduction is now a standard application of the Chinese remainder theorem. Given an Eulerian graph  $G = (V, E)$ , we can, w.l.o.g., assume that the degree of vertices of  $G$  is at least 4 (vertices of degree 2 can be removed by contracting edges). The number of Eulerian tours of a graph on  $n$  vertices is bounded by  $n^{n^2}$  (the number of pairings in a vertex of degree  $d$  is  $d!/(2^{d/2}(d/2)!) \leq n^n$ ).

We choose  $n^2$  primes  $p_1, \dots, p_{n^2} > n$  such that  $\prod_{i=1}^{n^2} p_i > n^{n^2}$  and each  $p_i$  is bounded by  $O(n^3)$  (see, e.g., [BS96], p.296). For each  $p_i$ , we construct graph  $G_i$  by replacing each vertex  $v$  of degree  $d > 4$  with  $Q$  gadget with  $d$  input and  $d$  output half-edges where the  $(2j - 1)$ -st and  $2j$ -th output half-edge are connected (for  $j = 1, \dots, d/2$ ), and the input half-edges are used to replace half-edges emanating from  $v$  (that is, they are connected to the input half-edges of other gadgets according to the edge incidence at  $v$ ). Note that  $G_i$  is a 4-regular graph. Since  $p_i = O(n^3)$ , the construction of  $G_i$  can be done in time polynomial in  $n$ . Having  $G_i$ , we make a query to the oracle and obtain the number  $T_i$  of Eulerian tours in  $G_i$ . Let  $T$  be the number of Eulerian tours in  $G$ . Then

$$T_i \equiv T \prod_{d=6}^n \left( \left( \frac{d}{2} \right)! 2^{d/2} R_d \right)^{n_d} \pmod{p_i}, \quad (6)$$

where  $n_d$  is the number of vertices of degree  $d$  in  $G$ .

Since  $T_i$  is of length polynomial in  $n$ , we can compute  $T_i \pmod{p_i}$  for each  $i$  and thus  $T \pmod{p_i}$  (since on the right hand side of (6)  $T$  is multiplied by a term that has an inverse modulo  $p_i$ ). By the Chinese remainder theorem, we can compute  $T$  in time polynomial in  $n$  (see, e.g., [BS96], p.106).  $\square$

### 3.3 #ET in 4-Regular Planar Graphs Is #P-Complete

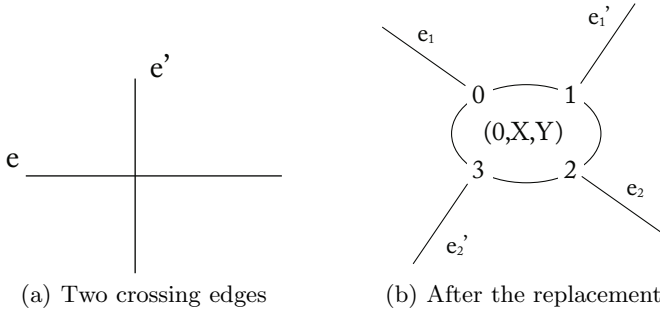
First, it's easy to see that #ET in 4-regular planar graphs is in #P. We will give a reduction from #ET in 4-regular graphs to #ET in 4-regular planar graphs.

**Theorem 2.** #ET in 4-regular graphs is polynomial time Turing reducible to #ET in 4-regular planar graphs.

*Proof.* Given a 4-regular graph  $G = (V, E)$ , we first draw  $G$  in the plane. We allow the edges to cross other edges, but i) edges do not cross vertices, ii) each crossing involves 2 edges. The embedding can be found in polynomial time.

Let  $p$  be an odd prime, we will construct a graph  $G_p$  from the embedded graph as follows. Let  $e, e'$  be two edges in  $G$  which cross in the plane as shown in Figure 4(a), we split  $e$  (and  $e'$ ) into two half-edges  $e_1, e_2$  ( $e'_1, e'_2$ , respectively). As illustrated in Figure 4(b), a  $(0, X, Y)$  node with parameters  $p$  and  $k = p$  is added, and  $e_1, e'_1, e_2, e'_2$  are connected to the half-edges 0,1,2,3 of the  $(0, X, Y)$  node, respectively.

Let  $G_p$  be the graph after replacing all crossings by  $(0, X, Y)$  nodes. We have that  $G_p$  is planar since  $(X, Y, Y)$  nodes and  $(0, X, Y)$  nodes are all planar. The construction can be done in time polynomial in  $p$  and the size of  $G$  (since the



**Fig. 4.** To replace a crossover point by a  $(0, X, Y)$  node with parameters  $p$  and  $k = p$

number of crossover points is at most  $O(|E|^2)$  and the size of each  $(0, X, Y)$  node is  $O(p^2)$ .

In the reduction, we choose  $n = |V|$  primes  $p_1, p_2, \dots, p_n$  such that  $p_i = O(n^2)$  for  $i \in [n]$  and  $\prod_{i=1}^n p_i \geq 3^n$ , where  $3^n$  is an upper bound for the number of Eulerian tours in  $G$  (the number of pairings in each vertex is 3). For each  $p_i$ , we construct a graph  $G_{p_i}$  from the embedded graph as described above with  $p = p_i$ . Let  $T$  be the number of Eulerian tours in  $G$  and  $T_i$  be the number of Eulerian tours in  $G_{p_i}$ , we have

$$T \equiv T_i \pmod{p_i}. \quad (7)$$

Equation (7) follows from the fact that the number of Eulerian tours in which the set of routes within any  $(0, X, Y)$  node is not of type  $\{\{0, 2\}, \{1, 3\}\}$  is zero (modulo  $p_i$ ) (since in (2) we have  $A \equiv 1 \pmod{p_i}$  and in (3) we have  $B \equiv 0 \pmod{p_i}$ ). We can make a query to the oracle to obtain the number  $T_i$ . By the Chinese remainder theorem, we can compute  $T$  in time polynomial in  $n$ .  $\square$

### 3.4 #A-TRAILS in 4-Regular Graphs with Rotational Embedding Schemes Is #P-Complete

In this section, we consider #A-TRAILS in graphs with rotational embedding schemes (maps). We prove that #A-TRAILS in 4-regular maps is #P-complete by a simple reduction from #ET in 4-regular graphs.

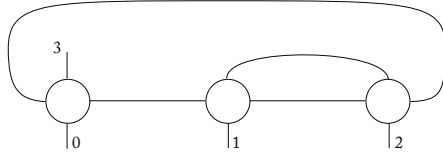
First, it's not hard to verify that #A-TRAILS in 4-regular maps is in #P.

**Theorem 3.** #ET in 4-regular graphs is polynomial time Turing reducible to #A-TRAILS in 4-regular maps.

*Proof.* Given a 4-regular graph  $G = (V, E)$ , for each vertex  $v$  of  $G$ , we use the gadget shown in Figure 5 to replace  $v$ .

The gadget consists of three vertices which are represented by circles in Figure 5. The labels 0, 1, 2 and 3 are the four half-edges which are used to replace half-edges emanating from  $v$ . The cyclic ordering of the 4 (half-)edges





**Fig. 5.** Gadget simulating vertex of degree 4

incident to each circle is given by the clockwise order, as shown in Figure 5. There are three types of valid sets of routes inside the gadget,  $VR(\{0, 1\}, \{2, 3\})$ ,  $VR(\{0, 2\}, \{1, 3\})$  and  $VR(\{0, 3\}, \{1, 2\})$ . By enumeration, we have the size of each of the three sets is 2.

Let  $G'$  be the 4-regular map obtained by replacing each vertex  $v$  by the gadget. Let  $T$  be the number of Eulerian tours in  $G$ , we have the number of A-trails in  $G'$  is  $2^{|V|}T$ .  $\square$

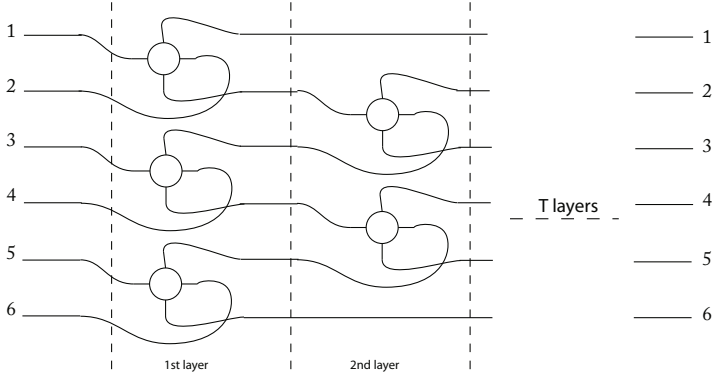
Note that Kotzig [Kot68] gave a one-to-one correspondence between the A-trails in any 4-regular plane graph  $G$  (the embedding in the plane gives the rotational embedding scheme) and the spanning trees in a plane graph  $G'$ , where  $G$  is the medial graph of  $G'$ . By the Kirchhoff's theorem (c.f. [Jer03]), the number of spanning trees of any graph can be computed in polynomial time. Thus #A-TRAILS in 4-regular plane graphs can be computed in polynomial time.

## 4 The Complexity of Approximate Counting

In this section, we show that #ET in general graphs is AP-reducible to #A-TRAILS in 4-regular maps. AP-reductions were introduced by Dyer, Goldberg, Greenhill and Jerrum [DGGJ04] for the purpose of comparing the complexity of two counting problems in terms of approximation (given two counting problems  $f, g$ , if  $f$  is AP-reducible to  $g$  and there is an FPRAS for  $g$ , then there is also an FPRAS for  $f$ ).

In the AP-reduction from #ET to #A-TRAILS in 4-regular maps, we use the idea of simulating the pairings in a vertex by a gadget as what we did in the construction of the  $Q$  gadget. The difference is that the new gadget works in an approximate way, that is, instead of having the number of valid sets of routes to be the same for each of the types, the numbers can be different but within a small multiplicative factor. The analysis of the gadget uses a fast mixing result for a card shuffling problem.

We use the gadget illustrated in Figure 6. The circles represent the vertices in the map. Let  $d$  be an even number. The gadget has  $d$  input half-edges on left and  $d$  output half-edges (Figure 6 demonstrates the case of  $d = 6$ ). There are  $T$  layers in the gadget which are numbered from 1 to  $T$  from left to right. In an odd layer  $t$ , the  $(2i - 1)$ -st and the  $2i$ -th output half-edges of layer  $t - 1$  are connected to a vertex of degree 4, for  $i \in [d/2]$ . In an even layer  $t$ , the  $2i$ -th and



**Fig. 6.** Construction of the gadget for a vertex of degree 6

the  $(2i + 1)$ -st output half-edges of layer  $t - 1$  are connected to a vertex of degree 4, for  $i \in [d/2 - 1]$ . In Figure 6, we illustrate the first two layers each of which is in two consecutive vertical dashed lines. The cyclic ordering of each vertex is given by the clockwise ordering (in the drawing in Figure 6), and so we have that the two half-edges in each vertex which are connected to half-edges of the previous layer are not cyclicly-adjacent.

Note that a valid route in the gadget always connects an input half-edge to an output half-edge. Thus a valid set of routes always realizes some permutation  $\sigma$  connecting input half-edge  $i$  to output half-edge  $\sigma(i)$ .

In order to prove that  $|VR(\sigma)|$  is almost the same for each permutation  $\sigma \in S_d$ , we show that for  $T = \Theta(d^2 \log d \log(d!/\varepsilon))$  we have  $|VR(\sigma)| / \sum_{\varrho \in S_d} |VR(\varrho)| \in [(1 - \varepsilon)/d!, (1 + \varepsilon)/d!]$  for each permutation  $\sigma \in S_d$ . The gadget can be interpreted as a process of a Markov chain for shuffling  $d$  cards. The simplest such chain proceeds by applying adjacent transpositions. The states of the chain are all the permutations in  $S_d$ . In each time step, let  $\sigma \in S_d$  be the current state, we choose  $i \in \{1, \dots, d - 1\}$  uniformly at random, and then switch  $\sigma(i)$  and  $\sigma(i + 1)$  with probability  $1/2$  and stay the same with probability  $1/2$ . For our gadget, it can be viewed as an even/odd sweeping Markov chain on  $d$  cards [Wil04]. The ratio  $|VR(\sigma)| / \sum_{\varrho \in S_d} |VR(\varrho)|$  is exactly the probability of being  $\sigma$  at time  $T$  when the initial state of the even/odd sweeping Markov chain is the identity permutation. By the analysis in [Wil04], we can relate  $T$  with the ratio as follows.

**Lemma 4** ([Wil04]). *Let  $T$  be the number of layers of the gadget with  $d$  input half-edges and  $d$  output half-edges as shown in Figure 6, and let  $\mu, \lambda$  be two distributions on  $S_d$  such that  $\mu(\sigma) = |VR(\sigma)| / \sum_{\varrho \in S_d} |VR(\varrho)|$  and  $\lambda(\sigma) = 1/d!$  ( $\lambda$  is the uniform distribution on  $S_d$ ). For*

$$T = O(d^2 \log d \log(d!/\varepsilon)),$$

*then  $\|\mu - \lambda\|_{TV} \leq \varepsilon/d!$ , and thus  $(1 - \varepsilon)/d! \leq \mu(\sigma) \leq (1 + \varepsilon)/d!$ .*

**Theorem 4.** *If there is an FPRAS for #A-TRAILS in 4-regular maps, then we have an FPRAS for #ET in general graphs.*

*Proof.* Given an Eulerian graph  $G = (V, E)$  and an error parameter  $\varepsilon > 0$ , we can, w.l.o.g., assume that the degree of vertices of  $G$  is at least 4 (vertices of degree 2 can be removed by contracting edges). We construct graph  $G'$  by replacing each vertex  $v$  of degree  $d > 2$  with a gadget with  $d$  input half-edges,  $d$  output half-edges and  $T_d = \Theta(d^2 \log d \log(4d!n/\varepsilon))$  layers where the  $(2i - 1)$ -st and  $2i$ -th output half-edge are connected (for  $1 \leq i \leq d/2$ ), and the input half-edges are used to replace half-edges emanating from  $v$  (that is, they are connected to the input half-edges of other gadgets according to the edge incidence at  $v$ ). We have that  $G'$  has  $O(n^2 T_n) = O(n^4 \log n(n \log n + \log(1/\varepsilon)))$  vertices and can be constructed in time polynomial in  $n$  and  $1/\varepsilon$ .

Let  $\mathcal{A}$  be an FPRAS for #A-TRAILS in 4-regular maps by the assumption of the theorem, we run  $\mathcal{A}$  on  $G'$  with error parameter  $\varepsilon/2$ . Let  $\mathcal{A}(G', \varepsilon/2)$  be the output of  $\mathcal{A}$  and  $N_A$  be the number of A-trails in  $G'$ , we have  $\mathcal{A}(G', \varepsilon/2) \in [e^{-\varepsilon/2} N_A, e^{\varepsilon/2} N_A]$  with probability at least  $2/3$ . This process can be done in time polynomial in the size of  $G'$  and  $1/\varepsilon$ , which is polynomial in  $n$  and  $1/\varepsilon$ .

Let  $D_d$  be the number of vertices in the gadget of  $d$  input half-edges and  $d$  output half-edges, and let  $R_d = 2^{D_d} 2^{d/2} (d/2)!/d!$  and  $R = \prod_{d=4}^n R_d^{n_d}$  where  $n_d$  is the number of vertices of degree  $d$  in  $G$ . Our algorithm  $\mathcal{B}$  will output

$$\mathcal{B}(G, \varepsilon) = \mathcal{A}(G', \varepsilon/2)/R. \quad (8)$$

We next prove that  $\mathcal{B}$  is an FPRAS for #ET in general graphs. For every Eulerian tour in  $G$ , the type of the pairing in each vertex in  $G$  is fixed. Note that each pairing corresponds to  $(d/2)!2^{d/2}$  permutations in a gadget with  $d$  input half-edges and  $d$  output half-edges. By Lemma 4, we have  $(1 - \varepsilon/(4n))2^{D_d}/d! \leq |VR(\sigma)| \leq (1 + \varepsilon/(4n))2^{D_d}/d!$  for each  $\sigma \in S_d$  where  $VR(\sigma)$  is counted in a gadget with  $d$  input half-edges and  $d$  output half-edges. Thus, the number of A-trails in  $G'$  which correspond to the same Eulerian tour in  $G$  is in  $[(1 - \varepsilon/(4n))^n R, (1 + \varepsilon/(4n))^n R]$ . Let  $N_E$  be the number of Eulerian tours in  $G$ , we have  $N_A \in [(1 - \varepsilon/(4n))^n R N_E, (1 + \varepsilon/(4n))^n R N_E]$ , and thus for  $\varepsilon \leq 2n$ ,  $N_A/R \in [e^{-\varepsilon/2} N_E, e^{\varepsilon/4} N_E]$  (the case when  $\varepsilon > 2n$  is trivial,  $\mathcal{B}$  can just output  $3^n$ ). Since  $\mathcal{A}(G', \varepsilon/2) \in [e^{-\varepsilon/2} N_A, e^{\varepsilon/2} N_A]$  with probability at least  $2/3$ , then by (8), we have  $\mathcal{B}(G, \varepsilon) \in [e^{-\varepsilon} N_E, e^{\varepsilon} N_E]$  with probability at least  $2/3$ . This completes the proof.  $\square$

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